Usable Crypto: Introducing miniLock

Nadim Kobeissi HOPE X, NYC, 2014





"Browsers are an environment that is hostile to cryptography"

- Malleability of the JavaScript runtime.
- The lack of low-level (system-level) programming access.
- DOM-style vulnerabilities (XSS)



"Code Delivery is a Chicken-Egg Problem"

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* What prevents web app code from being intercepted and modified by a "man in the middle"?

Fine, why not use SSL?



Quote from popular anti-JS crypto article

"You can [use SSL]. It's harder than it sounds, but you [can] safely transmit Javascript crypto to a browser using SSL. The problem is, having established a secure channel with SSL, you no longer need Javascript cryptography; you have "real" cryptography."

> What the author ignores: Unlike SSL, JavaScript cryptography protects data from server access. Also claims people are using JS Crypto to get around deploying SSL (???)

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No One Saw the Value of JavaScript Cryptography

- JS shifting from language of the web to language of everything
- Making JS crypto real means making crypto work in the world's most accessible language
- Huge privacy/security gains in a usable environment





Cryptocat: Encrypted Chat in the Browser

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- Open source app with over 200,000 users.
- Goal: make encrypted chat accessible, fun, and easy to use.
- Accessible no matter your background.



Basic Needs

- Secure cryptographic primitives (AES, SHA2, ECDH).
- Secure pseudorandom number generation.
- Secure code delivery.



Cryptographic Primitives

- Public key cryptography and digital signature algorithms depend on numbers much larger than 64-bit floating point.
- Big integers require a third-party library



Cryptographic Primitives

- Some algorithms are computationally expensive (RSA, Diffie-Hellman, DSA...)
- Web workers came to the rescue.



Cryptographic Primitives

- Multiple maturing
 libraries: SJCL, Crypto-JS, OpenPGPJS
- Crypto operations

 (bitwise, etc.) are
 surprisingly cleanly
 writeable in JavaScript.



One Round of AES in JavaScript and C

1 a2 = t0[a>>>24] ^ t1[b>>16 & 0xff] ^ t2[c>>8 & 0xff] ^ t3[d & 0xff] ^ key[4]; 2 b2 = t0[b>>>24] ^ t1[c>>16 & 0xff] ^ t2[d>>8 & 0xff] ^ t3[a & 0xff] ^ key[5]; 3 c2 = t0[c>>>24] ^ t1[d>>16 & 0xff] ^ t2[a>>8 & 0xff] ^ t3[b & 0xff] ^ key[6]; 4 d = t0[d>>>24] ^ t1[a>>16 & 0xff] ^ t2[b>>8 & 0xff] ^ t3[c & 0xff] ^ key[7];

1 t0 = Te0[s0 >> 24] ^ Te1[(s1 >> 16) & 0xff] ^ Te2[(s2 >> 8) & 0xff] ^ Te3[s3 & 0xff] ^ rk[4]; 2 t1 = Te0[s1 >> 24] ^ Te1[(s2 >> 16) & 0xff] ^ Te2[(s3 >> 8) & 0xff] ^ Te3[s0 & 0xff] ^ rk[5]; 3 t2 = Te0[s2 >> 24] ^ Te1[(s3 >> 16) & 0xff] ^ Te2[(s0 >> 8) & 0xff] ^ Te3[s1 & 0xff] ^ rk[6]; 4 t3 = Te0[s3 >> 24] ^ Te1[(s0 >> 16) & 0xff] ^ Te2[(s1 >> 8) & 0xff] ^ Te3[s2 & 0xff] ^ rk[7];

JavaScript Cryptography: Example of a Bug

90	90	<pre>multiParty.genPrivateKey = function() {</pre>
91		<pre>- rand = Cryptocat.randomString(32, 0, 0, 1);</pre>
92		<pre>- myPrivateKey = BigInt.str2bigInt(rand, 10);</pre>
	91	<pre>+ var rand = Cryptocat.randomString(64, 0, 0, 0, 1);</pre>
	92	<pre>+ myPrivateKey = BigInt.str2bigInt(rand, 16);</pre>
93	93	return myPrivateKey;

Thanks, weak typing.

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Secure Pseudorandomness

Math.random() relies on guessable entropy sources.

* window.crypto.ge
tRandomValues()
doesn't.







Code Delivery

- Browser apps. Chrome led

 a revolution, model
 adopted by Safari and
 Opera (Firefox lags behind.)
- Great features: Code signing, enhanced security, protection against XSS and in-line eval.



Code Delivery

- In some cases, signed
 browser apps have benefits
 over regular desktop apps!
- Strong separation from system level.
- Chrome: tab CPU sandboxing.





W3C Web Crypto API

- Native cryptographic primitives!
- A solid chance to mitigate side-channel attacks such as timing attacks.
- (Disclosure: I'm on that team)



W3C Web Crypto API

- Missing features:
- Modern algorithms (Curve25519)
- Key storage API



Sudden Acceptance of JS Cryptography

- Google publishes browser extension for GPG, own JS cryptography library
- Microsoft publishes Microsoft JavaScript Cryptography Library
- Thai Duong: "Why JavaScript cryptography is useful"





Stéphane Bortzmeyer @bortzmeyer

If you have issues with Javascript crypto, wait til you see C Web crypto API (native crypto for JavaScript) is at "I crypto.

2014-07-09, 5:35 AM



Remaining Problems Today: Weak Typing

 I want ECMAScript to have optional strong typing.

* var k = 5

- * var number(k) = 5
- Both would work, but the second one can throw a TypeError if you do k = 5 + 'Meow'



What I'm introducing today

- My new usable encryption software.
- Let's innovate on file encryption and sharing.
- Let's make it universally accessible.



File encryption software that does more with less.

Current status for file encryption

- Main contender: PGP
- Main use case: File attachments
- Classic public key management

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Datei Bearbeiten Ansicht Einfügen Format ?

----BEGIN PGP PUBLIC KEY BLOCK----Version: GnuPG v1.4.3-cvs (MingW32)

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Key management is awful

- * Generating keys
- Saving them on disk with passphrase
- * Sharing long public keys via email
- Storing other people's keys, authenticating via fingerprints, managing keys

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This is not convenient and we can do a lot better.

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Generating keys

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No key storage. No key files.

- miniLock asks you for a passphrase and uses it to generate your key identity.
- Enter that passphrase on any computer in the world, obtain the same persistent key identity.
- Nothing is ever stored anywhere.





miniLock uses miniLock IDs.

 miniLock IDs are shareable public keys that are 44 characters long.

* Here's mine: 9LbEGtYBXRf1s0bIyw qbhty7uA00TF0XdynV +fIJlDc=





User flow

- I enter my passphrase on a miniLock-capable computer and get my miniLock ID (always the same).
- I can send files to others using their ID.
- I can receive files sent to my ID.





This sucks less.

- No private key storage or management.
- No managing long key identities of others (miniLock IDs are tweetable!)
- miniLock IDs are so small that they act as their own fingerprint.




Nice features

- Easy to use interface
- Encrypt files for own use, decrypt later
- Runs on any computer





Nice features

* Send to multiple recipients (almost no performance decrease/file size increase)

 miniLock IDs of recipients are anonymized (even from the recipients)

* Fast!

Retains filename on decryption





Best of all

 Peer-reviewed design specification





Best of all

- Peer-reviewed design specification
- Fully audited (Thanks to Cure53 and OTF)





Unit Test Kit

- Simulates entire user flow with randomized use-cases
- Also can run
 independent user flow
 elements atomically



miniLock in your app!

- Highly portable
- Comes with full design documents/spec/tests/ reference
- Your app can be miniLock-ready





Everything will be released today

- Right after this talk
- AGPLv3 license
- But first...





How do the internals work?

- Reliance on elliptic curve cryptography (specifically, TweetNaCL)
- Mechanisms to evaluate strength of passphrases/ suggest strong pass phrases





* Scrypt.

TweetNaCL

 "World's first auditable highsecurity cryptographic library" — Daniel J. Bernstein

 Tiny, capable, easy to audit (fits in 100 tweets)

 Ported to JS by Dmitry Chestnykh





TweetNaCL

* Offers interface for:

Curve25519 (public key generation)

Xsalsa20 (Encryption)

* Poly1305 (Authentication)





Curve25519

32-byte private keys,
 32-byte public keys
 (tiny!)

* Extremely fast





Key derivation



Key derivation



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User enters passphrase

 Optimally, we want to map 32 bytes of entropy into the 32byte Curve25519 private key

Not practically feasible.





Key derivation



Entropy evaluation (100+ bits)



Suggest passphrase Derive Curve25519 key pair

Entropy Evaluation

- We measure user passphrase entropy (using zxcvbn)
- miniLock suggests "a long, unique phrase that makes sense only to you."





Entropy Evaluation

Less than 100-bit
 entropy pass phrases are
 not allowed (miniLock
 will refuse to open)

Instead, miniLock
 constructs a suggested
 passphrase





Key derivation



Entropy Evaluation

 miniLock ships with dictionary of 58,110 most used English words

7-word passphrase =
 58110⁷ ~= 2¹¹¹





Towards reliable storage of 56-bit secrets in human memory

Joseph Bonneau Princeton University Stuart Schechter Microsoft Research

Abstract

Challenging the conventional wisdom that users cannot remember cryptographically-strong secrets, we test the hypothesis that users can learn randomly-assigned 56bit codes (encoded as either 6 words or 12 characters) through *spaced repetition*. We asked remote research participants to perform a distractor task that required logging into a website 90 times, over up to two weeks, with a password of their choosing. After they entered their chosen password correctly we displayed a short code (4 letters or 2 words, 18.8 bits) that we required them to continue to be manufactured and deployed. But they are sufficiently pervasive that we must design our protocols around their limitations.)

-Kaufman, Perlman and Speciner, 2002 [60]

The dismissal of human memory by the security community reached the point of parody long ago. While assigning random passwords to users was considered standard as recently in the mid-1980s [29], the practice died out in the 90s [4] and NIST guidelines now presume all passwords are user-chosen [35]. Most banks have even given up on expecting customers to memorize random

100+ bits of entropy

- Sufficient our purposes
- We can also work on making it harder to map the keyspace





Key derivation



Entropy evaluation (100+ bits)

Suggest passphrase Derive Curve25519 key pair

Scrypt

Scrypt

- Provides "memory-hard" key derivation.
- First we derive a SHA-512 hash of the passphrase
- Hash goes through 2¹⁷
 rounds





Key derivation



Key derivation

- ★ scrypt(L) = 32 bytes →
 Curve25519 private key
- miniLock ID is Base64
 encoding of public key





- File is encrypted using a random unique symmetric key
- Symmetric key is encrypted asymmetrically once for each recipient and stored in header
- Both are authenticated encryption





Begin header

Sender ID

Recipient A: Header info

Recipient B: Header info

End header

Ciphertext

Begin header

Sender ID

Recipient A: Header info

Recipient B: Header info

End header

Ciphertext

64

File key File nonce File name

(Encrypted to
Recipient A's
public key)

B

File key File nonce File name

(Encrypted to
Recipient B's
public key)

Header recipients are anonymized

File key File nonce File name

(Encrypted to Recipient ?'s public key) File key File nonce File name

(Encrypted to
Recipient ?'s
public key)

Header recipients are anonymized

- Recipient attempts to decrypt every section of the header
- If they obtain an authenticated decryption, they know they are an intended recipient





General usage

- Share your miniLock IDs with friends
- Encrypt any files using friends' miniLock IDs
- Decrypt files sent to you
- Drag and drop simplicity





Demonstration





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Release schedule

- miniLock is audited, reviewed software: ready for use
- 2-week test period
 before "App Store"
 release





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Release schedule

- Will be released as a Chrome app
- Runs on Chrome OS, Windows, Mac, Linux
- 2-week test period before
 "App Store" release

Thank you!

 Get the code and documentation today

http://minilock.io



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